Arrangements of Pseudocircles: On Digons and Triangles

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Abstract

In this article, we study the cell-structure of simple arrangements of pairwise intersecting pseudocircles. The focus will be on two problems of Grünbaum (1972).

First, we discuss the maximum number of digons or touching points. Grünbaum conjectured that there are at most 2n-2 digon cells or equivalently at most 2n-2 touchings. Agarwal et al. (2004) verified the conjecture for cylindrical arrangements. We show that the conjecture holds for any arrangement which contains three pairwise touching pseudocircles. The proof makes use of the result for cylindrical arrangements. Moreover, we construct non-cylindrical arrangements which attain the maximum of 2n-2 touchings and have no triple of pairwise touching pseudocircles.

Second, we discuss the minimum number of triangular cells (triangles) in arrangements without digons and touchings. Grünbaum conjectured that such arrangements have 2n-4 triangles. Snoeyink and Hershberger (1991) established a lower bound of $\lceil \frac{4}{3}n \rceil$. Felsner and Scheucher (2017) disproved the conjecture and constructed a family of arrangements with only $\lceil \frac{16}{11}n \rceil$ triangles. We provide a construction which shows that $\lceil \frac{4}{3}n \rceil$ is the correct value.

Keywords and phrases arrangement of pseudocircles, touching, empty lense, cylindrical arrangement, arrangement of pseudoparabolas, Grünbaum's conjecture

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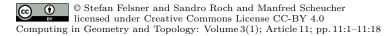
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1 Introduction

An arrangement \mathcal{A} of pairwise intersecting pseudocircles is a collection of $n(\mathcal{A})$ simple closed curves on the sphere or plane such that any two of the curves either touch in a single point or intersect in exactly two points where they cross. Throughout this article, we consider all arrangements to be simple, that is, no three pseudocircles meet in a common point. An arrangement \mathcal{A} partitions the plane into cells. A cell with exactly k crossings on its boundary is a k-cell, 2-cells are also called digons and 3-cells are triangles. The number of k-cells of an arrangement \mathcal{A} is denoted as $p_k(\mathcal{A})$.

The study of cells in arrangements started almost a century ago when Levi [11] showed that, in an arrangement of at least three pseudolines in the projective plane, every pseudoline





is incident to at least three triangles. In the 1970's, Grünbaum [10] studied arrangements of pseudolines and initiated the study of arrangements of pseudocircles.

1.1 Digons and touchings

Concerning digons in arrangements of pairwise intersecting pseudocircles, Grünbaum [10] presented a construction with 2n-2 digons (depicted in Figure 1) and conjectured that these arrangements have the maximum number of digons¹.

▶ Conjecture 1 (Grünbaum's digon conjecture [10, Conjecture 3.6]). Every simple arrangement A of n pairwise intersecting pseudocircles has at most 2n-2 digons, i.e., $p_2 \leq 2n-2$.

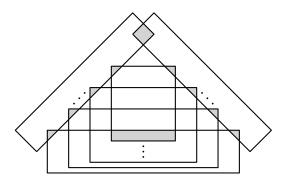


Figure 1 An arrangement of $n \ge 4$ pairwise intersecting pseudocircles with exactly 2n - 2 digons. Digons are highlighted gray (Example copied from Grünbaum [10, Figure 3.28]).

It was shown by Agarwal et al. [2] that Conjecture 1 holds for simple cylindrical arrangements.

An intersecting arrangement of pseudocircles is cylindrical if there is a pair of cells which are separated by each pseudocircle of the arrangement. An intersecting arrangement of pseudoparabolas is a collection of infinite x-monotone curves, called pseudoparabolas, where each pair of them either have a single touching or intersect in exactly two points where they cross. Every cylindrical arrangement of pseudocircles can be represented as an arrangement of pseudoparabolas and vice versa. From an arrangement of pseudoparabolas one can directly obtain a drawing of an arrangement of pseudocircles on the lateral surface of a cylinder so that the pseudocircles wrap around the cylinder. The two separating cells correspond to the top and the bottom of the cylinder. The top and bottom cell may have degree two, they account for the difference of two between the conjecture and the result for pseudo-parabolas.

More precisely Agarwal et al. [2] showed that the number of touchings in an intersecting arrangement of n pseudo-parabolas is at most 2n-4 [Theorem 2.4] and in an intersecting cylindrical arrangement it is at most 2n-3 [Corollary 2.12]. They extended this by showing that intersecting arrangements of pseudocircles the number of digons is at most linear in n [Theorem 2.13]. The proof of the last result is based on the fact that every arrangement of intersecting pseudocircles can be stabbed by constantly many points. That is, there exists an absolute constant k, called the $stabbing\ number^2$, such that for every arrangement of n pseudocircles in the plane there exists a set of k points with the property that each pseudocircle

¹ Originally the conjecture was stated as to include non-simple arrangements which are *non-trivial*, i.e., non-simple arrangements with at least 3 crossing points.

² In the literature, the stabbing number is also referred to as *piercing number* or *transversal number*.

contains at least one of the points in its interior region [2, Corollary 2.8]. Therefore, the arrangement can be decomposed into constantly many cylindrical subarrangements. The linear upper bound then follows from the fact that each pair of subarrangements contributes at most linearly many digons. Grünbaum's digon conjecture is known to hold for arrangements with up to 7 pseudocircles; see [8].

The conjecture has also been studied for proper circles. Alon et al. [3] proved that every arrangement of n pairwise intersecting circles contains at most 20n - 2 digons. For arrangements of pairwise intersecting unit circles, Pinchasi [12] proved that there are at most n + 3 digons. Very recently Ackerman et al. [1] verified Grünbaums digon conjecture for circles.

In this paper we show that Grünbaum's digon conjecture (Conjecture 1) holds for arrangements which contain three pseudocircles that pairwise form a digon. Before we state the result as a theorem, let us introduce some notation. For an arrangement \mathcal{A} of pseudocircles and any selection of its digons, we can perform a perturbation so that the selected digons become touching points. Figure 2 gives an illustration. It is therefore sufficient to find an upper bound on the number of touchings to prove Grünbaum's digon conjecture. We define the touching graph $T(\mathcal{A})$ to have the pseudocircles of \mathcal{A} as vertices, and two vertices form an edge if the two corresponding pseudocircles touch.

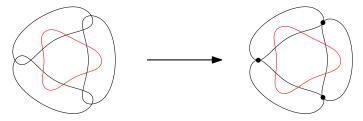


Figure 2 Contracting some digons to touchings.

▶ **Theorem 2.** Let A be a simple arrangement of n pairwise intersecting pseudocircles. If the touching graph T(A) contains a triangle, then there exist at most 2n-2 touchings, i.e., $p_2(A) \leq 2n-2$.

Theorem 2 in particular shows that Grünbaum's construction with 2n-2 touchings is maximal for arrangements with triangles in the touching graph. However, the maximum number of touchings in general arrangements remains unknown. In Section 3 we construct a family of arrangements of n pairwise intersecting pseudocircles which have exactly 2n-2 touchings and a triangle free touching graph. This family witnesses that the conjectured upper bound (Conjecture 1) can also be achieved in the cases not covered by Theorem 2.

▶ Proposition 3. For $n \in \{11, 14, 15\}$ and $n \ge 17$ there exists a simple arrangement \mathcal{A}_n of n pairwise intersecting pseudocircles with no triangle in the touching graph $T(\mathcal{A}_n)$ and with exactly $p_2(\mathcal{A}_n) = 2n - 2$ touchings.

1.2 Triangles in digon-free arrangements

In this context we assume that all arrangements are digon- and touching-free. It was shown by Levi [11] that every arrangement of n pseudolines in the projective plane contains at least n triangles. Since arrangements of pseudolines are in correspondence with arrangements

of great-pseudocircles (see e.g. [7, Section 4]), it follows directly that an arrangement of n great-pseudocircles contains at least 2n triangles, i.e., $p_3 \ge 2n$.

Grünbaum conjectured that every digon-free intersecting arrangement on n pseudocircles contains at least 2n-4 triangles [10, Conjecture 3.7]. Snoeyink and Hershberger [14] proved a sweeping lemma for arrangements of pseudocircles. Using this powerful tool, they concluded that in every digon-free intersecting arrangement every pseudocircle has two triangles on each of its two sides (interior and exterior). This immediately implies the lower bound $p_3(A) \ge 4n/3$; see Section 4.2 in [14].

An infinite family of digon-free arrangements of intersecting pseudocircles with $p_3 < \frac{16}{11}n$ was constructed in [8]. This family shows that Grünbaum's conjecture is wrong. With computer assistence [8], it was also verified that the lower bound $p_3 \ge 4n/3$ is tight for $6 \le n \le 14$. Here we show that the bound is tight for all $n \ge 6$:

▶ **Theorem 4.** For every $n \ge 6$, there exists a simple digon-free arrangement \mathcal{A}_n of n pairwise intersecting pseudocircles with $p_3(\mathcal{A}_n) = \lceil \frac{4}{3}n \rceil$ triangles. Moreover, these arrangements are cylindrical.

The construction that we use for proving Theorem 4 is based on replacing pseudocircles by bundles of pseudocircles. Starting from any digon-free arrangement, we can extend it to a counterexample of Grünbaum's triangle conjecture:

▶ Theorem 5. For $\varepsilon > 0$ fixed, every digon-free arrangement \mathcal{A} of n pairwise intersecting pseudocircles is contained as a subarrangement in a digon-free arrangement $\hat{\mathcal{A}}$ of \hat{n} pairwise intersecting pseudocircles with $p_3(\hat{\mathcal{A}}) < (\frac{4}{3} + \varepsilon)\hat{n}$.

For n=6 there is a unique intersecting digon-free arrangement which minimizes the number of triangles. This arrangement \mathcal{N}_6^{Δ} (see Figure 17) has been shown to be non-circularizable in [7], i.e., \mathcal{N}_6^{Δ} cannot be represented as an arrangement of proper circles. All counterexamples to Grünbaums conjecture presented in [8] as well as the arrangements constructed in the proof of Theorem 4 contain \mathcal{N}_6^{Δ} as a subarrangement. Hence, all these arrangements are non-circularizable. The following weakening of the original Grünbaum conjecture has been stated as Conjecture 2.2 in [8].

▶ Conjecture 6 (Weak Grünbaum triangle conjecture). Every simple digon-free arrangement A of n pairwise intersecting circles has at least 2n-4 triangles.

To prove the conjecture it would be enough to verify that every simple intersecting digon-free arrangement of n pseudocircles with less than 2n-4 triangles contains \mathcal{N}_6^{Δ} as a subarrangement. This, however, is wrong. There are counterexamples to Grünbaum's conjecture without \mathcal{N}_6^{Δ} as a subarrangement. In Subsection 4.2, we prove the following proposition and discuss additional constructions.

▶ Proposition 7. There is an infinite family of simple intersecting digon-free arrangements of n pseudocircles with $\lceil \frac{5}{3}n \rceil + 2$ triangles which have no subarrangement isomorphic to \mathcal{N}_6^{Δ} .

1.3 Related Work and Discussion

In the proof of Theorem 2, we make use of a triangle (K_3) in the touching graph to bound the number of digons in the arrangement. It would be interesting to know whether other subgraphs like C_4 or $K_{3,3}$ can also be used to bound the number of digons.

The focus of this article is on arrangements of pairwise intersecting pseudocircles. For the setting of arrangements, where pseudocircles do not necessarily pairwise intersect, a classical

construction of Erdős [5] gives arrangements of n unit circles with $\Omega(n^{1+c/\log\log n})$ touchings. An upper bound of $O(n^{3/2+\epsilon})$ on the number of digons in *circle* arrangements was shown by Aronov and Sharir [4]. The precise asymptotics, however, remain unknown. Moreover, we are not aware of an upper bound for *pseudocircles*.

ightharpoonup Problem 1. Determine the maximum number of touchings among all simple non-intersecting arrangements of n circles and pseudocircles, respectively.

Concerning intersecting arrangements with digons, the number of triangles behaves different than in digon-free arrangements. While the lower bound is $p_3 \geq 2n/3$, we know that in the range of $3 \leq n \leq 7$ the correct bound is $p_3 \geq n-1$, this was obtained using a computer-assisted exhaustive enumeration [8]. This motivates the conjecture:

▶ Conjecture 8 ([8, Conjecture 2.10]). Every simple arrangement of $n \ge 3$ pairwise intersecting pseudocircles has at least n-1 triangles, i.e., $p_3 \ge n-1$.

Dropping the condition that the arrangements are intersecting seems non-interesting at first: arrangements of pairwise non-intersecting circles have no triangles. If we add the condition that the intersection graph of the circles is connected we still have arrangements with a bipartite intersection graph where all faces are of even length, hence, there are no triangles. In the case where the intersection graph is connected and digons are forbidden triangles are unavoidable, in fact for all $n \geq 3$ there are arrangements in this class with only 8 triangles and this is the minimum.

Concerning the maximum number of triangles in intersecting arrangements, Felsner and Scheucher [8] have shown an upper bound $p_3 \leq \frac{4}{3} \binom{n}{2} + O(n)$ which is optimal up to a linear error term. In fact, while $\frac{4}{3} \binom{n}{2}$ is an upper bound for arrangements of great-pseudocircles, we found an intersecting arrangement (n=7) with no digons, no touchings, and $29 = \frac{4}{3} \binom{7}{2} + 1$ triangles. We are not aware of an infinite family of such arrangements.

▶ **Problem 2.** Determine the maximum number of triangles of simple arrangements of n pairwise intersecting pseudocircles.

2 Proof of Theorem 2

▶ **Theorem 2.** Let A be a simple arrangement of n pairwise intersecting pseudocircles. If the touching graph T(A) contains a triangle, then there exist at most 2n-2 touchings, i.e., $p_2(A) \leq 2n-2$.

Proof. Since the touching graph T(A) contains a triangle, there are three pseudocircles in A that pairwise touch. Let K be the subarrangement induced by these three pseudocircles and let Δ and Δ' denote the two open triangle cells in K. We label the three touching points, which are also the vertices of Δ and Δ' , as a, b, c. Furthermore, we label the three boundary arcs of Δ (resp. Δ') as α, β, γ (resp. α', β', γ'), as shown in Figure 3a.

Assume that all digons in \mathcal{A} are contracted to touchings. In the following, the arrangement in Figure 5 will serve as a running example for \mathcal{A} . The intersection of a pseudocircle $C \in \mathcal{A} \setminus \mathcal{K}$ with $\Delta \cup \Delta'$ results in three connected segments, which we denote as the three *pc-arcs* of C; see Figures 3b and 3c. Note that two of the pc-arcs induced by C may share an endpoint if C forms a touching with one of the pseudocircles from \mathcal{K} ; in the example arrangement in Figure 5, this occurs 5 times on the boundary of Δ and 4 times on the boundary of Δ' .

Each pc-arc in \triangle connects two of α, β or γ while a pc-arc in \triangle' connects two of α', β' and γ' . Depending on the boundary arcs on which they start and end, they belong to one of

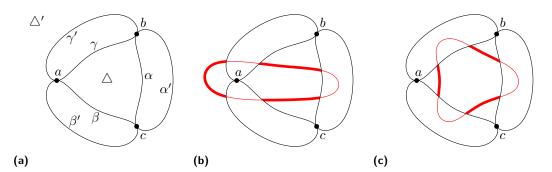


Figure 3 (a) An illustration of the subarrangement \mathcal{K} . (b) and (c), respectively, illustrate an additional pseudocircle C (red), the pc-arcs inside $\triangle \cup \triangle'$ are highlighted.

the types $\alpha\beta$, $\beta\gamma$, $\alpha'\gamma$, $\alpha'\beta'$, $\beta'\gamma'$ or $\alpha'\gamma'$. Figure 6 shows the regions of \triangle and \triangle' together with the pseudocircles passing through them; Figure 7 shows the same regions but the arcs are colored according to their type in blue $(\alpha\beta, \alpha'\beta')$, red $(\beta\gamma, \beta'\gamma')$ or blue $(\alpha\gamma, \alpha'\gamma')$.

 \triangleright Claim 9. If two pc-arcs inside \triangle (resp. \triangle') have a touching or cross twice, then they are of the same type.

Proof of Claim 9. We prove the claim for \triangle ; the argument for \triangle' is the same. Suppose towards a contradiction that two distinct pseudocircles C, C' from $A \setminus \mathcal{K}$ contain pc-arcs $A \subset C \cap \triangle$ and $A' \subset C' \cap \triangle$ of different types that have a touching or cross twice. For simplicity, consider only the arrangement induced by the five pseudocircles $\mathcal{K} \cup \{C, C'\}$. By symmetry we may assume that A is of type $\alpha\gamma$ and A' is of type $\alpha\beta$. We may further assume that A and A' have a touching, since otherwise, if they cross twice, they form a digon and we can contract it. This allows us to distinguish four cases which are depicted in Figure 4 (up to further possible contractions of digons formed between C and the pseudocircles of \mathcal{K}).

Case 1: C separates a from b and c.

Case 2: C separates b from a and c.

Case 3: C separates c from a and b.

Case 4: C does not separate a, b, c.

In the next paragraph we show that in neither case is it possible to extend the arc A' to a pseudocircle C' intersecting the three pseudocircles of K. This is a contradiction.

Extend A' starting from its endpoint on α . The only way to reach γ or γ' , avoiding an invalid, additional intersection with C, is via the pseudocircle $\beta \cup \beta'$. But the other endpoint of A' already lies on β , so either the pseudocircle extending A' has at least three intersections with $\beta \cup \beta'$ or it misses $\gamma \cup \gamma'$. Both are prohibited in an intersecting arrangement extending K.

This completes the proof of Claim 9.

 \triangle

Next we transform \mathcal{A} into another intersecting arrangement \mathcal{A}' by redrawing the pc-arcs within \triangle and \triangle' such that the pairwise intersections and touchings are preserved and all crossings and touchings of each arc type are concentrated in a narrow region as depicted in Figure 8. First we apply an appropriate homeomorphism on the drawing so that \triangle becomes a proper triangle (\triangle' will be treated in an analogous manner); see again Figure 6 and Figure 7. For the arc type $\alpha\beta$ we place a small rectangular region $R_{\alpha\beta}$ within \triangle that lies close to the vertex c. We now redraw all pc-arcs of type $\alpha\beta$ so that

- \blacksquare all crossings and touchings between pc-arcs of type $\alpha\beta$ lie inside $R_{\alpha\beta}$,
- every pc-arc of type $\alpha\beta$ intersects $R_{\alpha\beta}$ on opposite sites, and

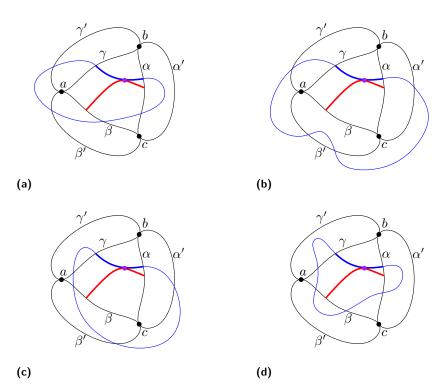


Figure 4 (a)–(d) illustrate Cases 1–4 from the proof of Claim 9. The pseudocircles C and C' are highlighted blue and red, respectively. The pc-arcs A and A' are emphasized.

for every pc-arc of type $\alpha\beta$, the removal of $R_{\alpha\beta}$ leaves two straight line segments which connect $R_{\alpha\beta}$ to α and β (i.e., the boundary segments of \triangle).

We proceed analogously for the arc types $\alpha\gamma$ and $\beta\gamma$. By Claim 9 touchings and double crossings only occur between pc-arcs of the same type and therefore lie in the rectangular regions. Since the rectangular regions are placed close enough to the vertices a,b,c of the triangle \triangle , no additional intersections or touching points are introduced and we obtain an arrangement \mathcal{A}' of pseudocircles with the same intersections and touchings as \mathcal{A} . The combinatorics of the resulting arrangement \mathcal{A}' may however differ from \mathcal{A} since the transformation typically changes the intersection orders of the pseudocircles. We conclude:

▶ Observation 10. The transformation preserves the incidence relation between any pair of pc-arcs, that is, two pc-arcs in \mathcal{A} are disjoint/cross in one point/cross in two points/touch if and only if the two corresponding pc-arcs in \mathcal{A}' are disjoint/cross in one point/cross in two points/touch.

This implies that \mathcal{A}' is indeed again an arrangement of $n(\mathcal{A}') = n(\mathcal{A})$ pairwise intersecting pseudocircles with identical touching graph $T(\mathcal{A}') = T(\mathcal{A})$. In particular, the number of touchings is preserved.

 \triangleright Claim 11. The arrangement induced by $\mathcal{A}' \setminus \mathcal{K}$ is cylindrical.

Proof of Claim 11. For each pseudocircle $C \in \mathcal{A}' \setminus \mathcal{K}$, the intersection

$$C\cap(\triangle\cup\triangle')=(C\cap\triangle)\cup(C\cap\triangle')$$

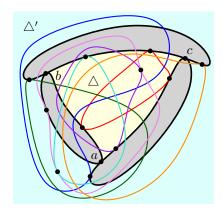


Figure 5 Original arrangement A

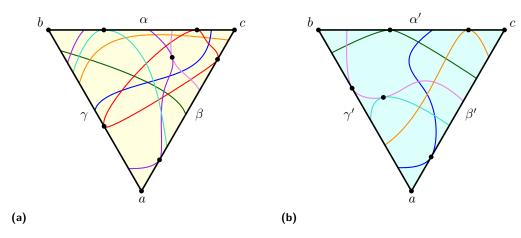


Figure 6 Area \triangle (a) and area \triangle' (b), colors as in \mathcal{A} .

consists of three pc-arcs, and each of these three pc-arcs is of a different type. The first arc is of type $\alpha\beta$ or $\alpha'\beta'$ (depending on whether it is inside Δ or Δ'), the second is of type $\beta\gamma$ or $\beta'\gamma'$, and the third is of type $\alpha\gamma$ or $\alpha'\gamma'$.

Now we redraw \mathcal{A}' on a cylinder as illustrated in Figure 9. Since all crossings and touchings of the arc type are within a small region, all pseudocircles from $\mathcal{A}' \setminus \mathcal{K}$ wrap around the cylinder, and hence the arrangement induced by $\mathcal{A}' \setminus \mathcal{K}$ is cylindrical.

This completes the proof of Claim 11.
$$\triangle$$

Next we replace the three pseudocircles of \mathcal{K} by six pseudocircles as illustrated in the left part of Figure 10, so that the resulting arrangement \mathcal{A}'' is cylindrical. Each of the three touching points a,b,c in \mathcal{K} is replaced by two new touching points and altogether we obtain touchings a',a'',b',b'',c',c''. Hence, when transforming \mathcal{A} into \mathcal{A}'' , the number of pseudocircles is increased by 3 and the number of touchings is also increased by 3.

Agarwal et al. [2] proved the $p_2 \leq 2n-3$ upper bound on the number of touchings in cylindrical arrangements of n pairwise intersecting pseudocircles by bounding the number of touchings in an arrangement of pairwise intersecting pseudoparabolas. They show that their touching graph is planar and bipartite [2, Theorem 2.4], hence, it has at most 2n-4 edges. The difference between 2n-4 and 2n-3 comes from the fact that the upper or the lower face in the pseudoparabola drawing of a pseudocircle arrangement \mathcal{A} can be a digon of \mathcal{A} .

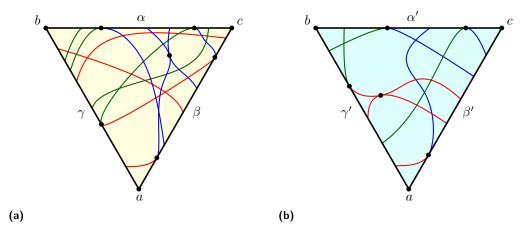


Figure 7 Area \triangle (a) and area \triangle' (b), colors according to arc-type.

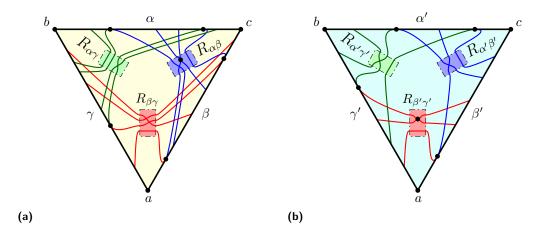


Figure 8 Area \triangle (a) and area \triangle' (b). Concentrate all crossings and touchings of one arc type in a narrow region. The narrow regions are indicated by dashed rectangles.

The drawing of \mathcal{A}'' in Figure 10 can be seen as an intersecting arrangement of pseudoparabolas. The complexity of the upper and of the lower face is three, hence, the arrangement has at most $2n(\mathcal{A}'') - 4$ touchings.

We review the ideas of the proof of Agarwal et al. [2] to verify the following claim.

 \triangleright Claim 12. $T(\mathcal{A}'')$ is planar, bipartite, and has at most $2n(\mathcal{A}'') - 5$ edges.

Proof of Claim 12. Label the pseudoparabolas P_1, \ldots, P_n such that the starting segments are ordered from top to bottom. In the touching graph $T(\mathcal{A}'')$, we label the corresponding vertices as $1, \ldots, n$.

Bipartiteness: The bipartition comes from the fact that the digons incident to a fixed pseudoparabola P_j are either all from below or all from above. Suppose that a pseudoparabola P_j has a touching from above with P_i and from below with P_k . It follows that P_i is above P_j everywhere and P_k is below P_j everywhere. Hence, P_i and P_k are separated by P_j and cannot intersect – this contradicts the assumption that the pseudocircles are pairwise intersecting.

We now further observe that the uppermost pseudoparabola P_1 and the lowermost pseudoparabola P_n belong to distinct parts of the bipartition, because P_1 has all touchings

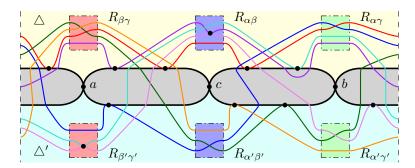


Figure 9 A cylindrical drawing of $\mathcal{A}' \setminus \mathcal{K}$

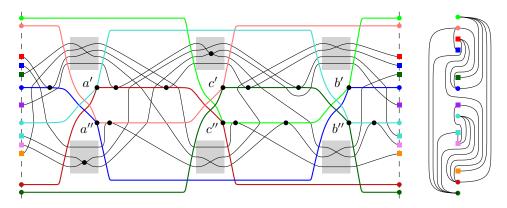


Figure 10 Replace each of the three pseudocircles of \mathcal{K} by two new pseudocircles so that the entire arrangement is now cylindrical. The pseudocircle from Figure 9 that contains the points a and c (resp. points c and b / points b and a) is replaced by a new dark red and a new bright red (resp. dark green and bright green / dark blue and turqoise) pseudocircle. The right part shows the corresponding touching graph.

below (i.e. with parabolas of greater index); P_n has all touchings above (i.e. with parabolas of smaller index). Hence, the touching graph remains bipartite after adding the edge $\{1, n\}$.

Planarity: For the planarity of $T(\mathcal{A}'')$, Agarwal et al. [2] create a particular drawing: The vertices are drawn on a vertical line and each edge $e = \{u, v\}$ is drawn as an y-monotone curve according to the following drawing rule: For each w with u < w < v, we route e to the left of w if the pseudoparabola P_w intersects P_u before P_v , otherwise we draw the edge e right of w. It is then shown that in the so-obtained drawing \mathcal{D} , each pair of independent edges has an even number of intersections. The right part of Figure 10 shows such a drawing of the corresponding touching graph. The Hanani–Tutte theorem (cf. Section 3 in [13]) implies that $T(\mathcal{A}'')$ is planar.

Notice that $\{1, n\}$ is not an edge in $T(\mathcal{A}'')$, since by construction, the lowermost and uppermost pseudocircles do not touch. We further observe that, since all edges in \mathcal{D} are drawn as y-monotone curves, the entire drawing lies in a box which is bounded from above by vertex 1 and from below by vertex n. Hence, we can draw an additional edge from 1 to n which is routed entirely outside of the box and does not intersect any other edge. Again, by the Hanani–Tutte theorem, we have planarity. Since any planar bipartite graph on n vertices has at most 2n-4 edges, we conclude that $T(\mathcal{A}'')$ has at most 2n-5 edges.

We are now ready to finalize the proof of Theorem 2. From Claim 12 and $n(\mathcal{A}'') = n + 3$ we get $p_2(\mathcal{A}'') \leq 2(n+3) - 5$. Since $p_2(\mathcal{A}'') = p_2(\mathcal{A}) + 3$ this implies $p_2(\mathcal{A}) \leq 2n - 2$, which is the desired bound.

3 Proof of Proposition 3

The proof of Proposition 3 is based on the *blossom operation*, which allows to dissolve triangles in the touching graph. We will apply the blossom operation to arrangements whose touching graphs are wheel graphs to obtain arrangements with the desired properties.

Let \mathcal{A} be an arrangement of pairwise intersecting pseudocircles, let v be a pseudocircle in \mathcal{A} , and let w_1, \ldots, w_d be the pseudocircles in \mathcal{A} which form touchings with v in this particular circular order along v. Since \mathcal{A} is intersecting, all the touchings are on the same side of v. As illustrated in Figure 11, the blossom operation relaxes the touchings between v and w_1, \ldots, w_d to digons and inserts d new pseudocircles v'_1, \ldots, v'_d inside and very close to v so that

- v'_1, \ldots, v'_d form a cylindrical arrangement,
- v touches v'_1, \ldots, v'_d , and
- w_i touches v'_{i-1} and v'_i (indices modulo d).

Since the new pseudocircles v'_1, \ldots, v'_d are added in an ε -small area close to v, it is ensured that each v'_i intersects all other pseudocircles. Hence, the obtained arrangement is again an arrangement of pairwise intersecting pseudocircles.

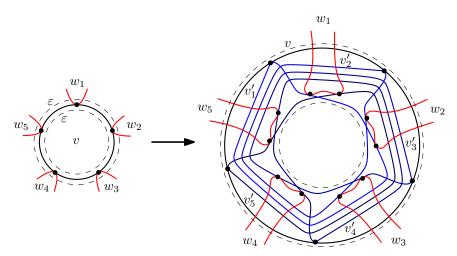


Figure 11 An illustration of the blossom operation applied on the pseudocircle v of an arrangement.

Figure 12 shows the effect of the blossom operation on the touching graph. Note that in these graph drawings the circular orders of the edges incident to a vertex coincide with the orders in which the touchings appear on the corresponding pseudocircle.

The blossom operation increases the number n(A) of pseudocircles of arrangement A by d and the number $p_2(A)$ of touchings by 2d. Hence, when applied to an arrangement A with exactly $p_2(A) = 2n(A) - 2$ touchings, the blossom operation again yields an arrangement A' with $p_2(A') = 2n(A') - 2$ touchings.

The blossom operation can be used to eliminate triangles in the touching graph. Assume pseudocircles w_i and w_j touch, hence v, w_i, w_j form a triangle in the touching graph. Then the blossom operation on v destroys this triangle without creating a new one if and only if,

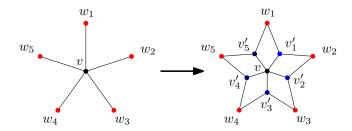


Figure 12 Blossom operation applied on v: Modification of the touching graph.

along the pseudocircle v, the two touchings with w_i and w_j are not consecutive. In Figure 12 a triangle $\{v, w_1, w_2\}$ would result in the new triangle $\{v'_1, w_1, w_2\}$, while a triangle $\{v, w_1, w_3\}$ would be eliminated without replacement.

Using the blossom operation, we are now able to prove Proposition 3.

▶ Proposition 3. For $n \in \{11, 14, 15\}$ and $n \ge 17$ there exists a simple arrangement \mathcal{A}_n of n pairwise intersecting pseudocircles with no triangle in the touching graph $T(\mathcal{A}_n)$ and with exactly $p_2(\mathcal{A}_n) = 2n - 2$ touchings.

Proof. Let $n' \geq 11$ be an integer with $n' \equiv 3 \pmod{4}$. Then $n = \frac{n'+1}{2}$ is an even integer with $n \geq 6$. As illustrated in Figure 13a and Figure 13b, we can construct an arrangement \mathcal{A} of n pseudocircles with $p_2 = 2n - 2$ touchings such that the touching graph $T(\mathcal{A})$ is the wheel graph W_n .

In this construction the *central* pseudocircle v has a touching with each of the pseudocircles w_1, \ldots, w_{n-1} and each w_i touches v, $w_{i+n/2}$, and $w_{i-n/2}$ (indices modulo n-1); see Figure 13c.

All triangles in $T(\mathcal{A})$ contain the central vertex v and for each such triangle $\{v, w_i, w_j\}$, the touchings of the pseudocircles w_i and w_j with the pseudocircle v are not consecutive on v. Therefore, applying the blossom operation to v eliminates all triangles and the resulting arrangement \mathcal{A}' of n' = 2n - 1 pairwise intersecting pseudocircles has $p_2(\mathcal{A}') = 2n' - 2$ touchings and a triangle-free touching graph $T(\mathcal{A}')$; see Figure 13d. This completes the argument for $n' \geq 11$ with $n \equiv 3 \pmod{4}$.

To give a construction for n'' = 14 and for all integers $n'' \ge 17$, note that the blossom operation can be applied to pseudocircles with exactly three touchings. The constructed examples with $n \equiv 3 \pmod{4}$ have pseudocircles with three touchings and the blossom operation applied to such a pseudocircle preserves the property.

Since n'' = 14 and every integer $n'' \ge 17$ can be written as n' + 3k with $n' \in \{11, 15, 19\}$ and $k \in \mathbb{N} \cup \{0\}$, we obtain arrangements \mathcal{A}'' of n'' pseudocircles with $p_2(\mathcal{A}'') = 2n'' - 2$ touchings. This completes the proof of Proposition 3.

4 Digon-free arrangements with few triangles

4.1 Proof of Theorem 4 and Theorem 5

The proofs of Theorem 4 and Theorem 5 are both based on replacing pseudocircles by canonical bundles of 4 pseudocircles, as shown in Figure 14. Like in the blossom operation, the new pseudocircles are placed within an ε -small area around the replaced pseudocircle so that the intersecting property of the pseudocircles is being preserved. We call such an operation a bundle replacement and aim for performing them on pseudocircles in order to destroy some of their incident triangles.

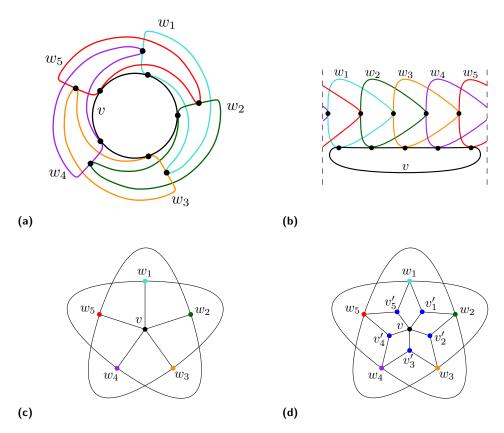


Figure 13 (a) An arrangement \mathcal{A} of 6 pseudocircles, (b) its cylindrical representation, (c) its touching graph $T(\mathcal{A})$, and (d) the touching graph $T(\mathcal{A}')$ after applying the blossom operation to v.

The following fact is a direct consequence of a sweepability statement by Snoeyink and Hershberger [14].

- ▶ Proposition 13 ([14, Lemma 4.1]). In a digon-free arrangement of pairwise intersecting pseudocircles, each pseudocircle is incident to at least two triangles to the inside and two triangles to the outside.
- ▶ Theorem 5. For $\varepsilon > 0$ fixed, every digon-free arrangement \mathcal{A} of n pairwise intersecting pseudocircles is contained as a subarrangement in a digon-free arrangement $\hat{\mathcal{A}}$ of \hat{n} pairwise intersecting pseudocircles with $p_3(\hat{\mathcal{A}}) < (\frac{4}{3} + \varepsilon)\hat{n}$.

Proof. Let \mathcal{A} be a digon-free arrangement of intersecting pseudocircles. Select a pseudocircle p in \mathcal{A} . Figure 15 illustrates a situation which might have been obtained via a bundle replacement on p. Note that a bundle replacement leads to a new arrangement \mathcal{A}' which is again digon-free and contains \mathcal{A} as a subarrangement.

We can think of the new bundle as being composed of four sections that are delimited by the four twists starting and ending in the purple crossings; colors refer to Figure 15. For a precise description we define a *twist* in a bundle as a sequence of consecutive crossings which make the outermost pseudocircle of the bundle the innermost. In all of our figures we keep the crossings of a twist close together.

Proposition 13 guarantees that p is incident to at least 4 triangles. Observe that the four twists in the bundle can always be distributed in such a way that each of these 4 triangles becomes incident to one of the twists, hence, the triangles are turned into quadrangles (green

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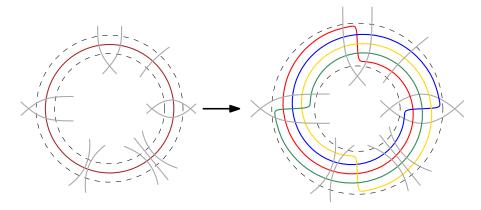


Figure 14 Bundle replacement: a pseudocircle is replaced by a bundle of 4 pseudocircles.

cells). On the other hand, in each of the 8 red areas, independent of the number of crossing pseudocircles (gray), exactly one new triangle is created.

In total, a careful bundle replacement on p leads to a digon-free arrangement \mathcal{A}' with $p_3(\mathcal{A}') \leq p_3(\mathcal{A}) + 4$ which contains \mathcal{A} as a subarrangement. This procedure can be iterated. Each iteration increases the number of pseudocircles by 3 and the number of triangles by at most 4.

Let t be the number of triangles of the initial arrangement \mathcal{A} . If $m \geq t/3\varepsilon$, then the arrangement $\hat{\mathcal{A}}$ obtained through a sequence of m bundle replacements has the claimed property: $p_3(\hat{\mathcal{A}}) \leq t + 4m \leq (4/3 + \varepsilon)3m < (4/3 + \varepsilon)\hat{n}$.

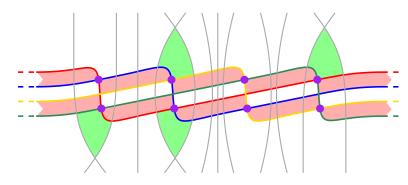
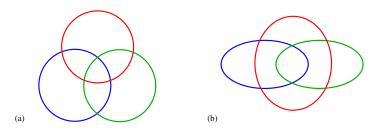


Figure 15 Situation obtained by a bundle replacement.



■ Figure 16 The two intersecting simple arrangements of three pseudocircles. (a) Krupp, (b) NonKrupp.

▶ **Theorem 4.** For every $n \ge 6$, there exists a simple digon-free arrangement \mathcal{A}_n of n pairwise intersecting pseudocircles with $p_3(\mathcal{A}_n) = \lceil \frac{4}{3}n \rceil$ triangles. Moreover, these arrangements are cylindrical.

Proof. Figure 16 shows the two intersecting simple arrangements of three pseudocircles. The Krupp is digon free and has 8 triangle cells. Note that the Krupp can be obtained from a single isolated circle in a bundle replacement step with a bundle of size three. Replacing any of the circles of the Krupp with a bundle of size four we can convert all the original triangles to 4-gons while generating eight new triangles. The result is the arrangement \mathcal{N}_6^{Δ} from Figure 17 with 6 pseudocircles and eight triangles. Starting from \mathcal{N}_6^{Δ} , we can iterate the bundle replacement with bundles of size four; this yields a family of arrangements \mathcal{A}_{3k} with n=3k pseudocircles and $4k=\frac{4}{3}n$ triangles.

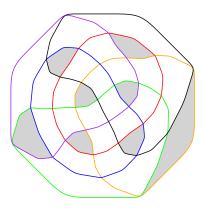


Figure 17 The non-circularizable digon-free intersecting arrangement \mathcal{N}_6^{Δ} .

For values of n which are not divisible by 3, we can take the arrangement \mathcal{A}_{3k} with $k = \lfloor \frac{n}{3} \rfloor$ and apply a bundle replacement step with a bundle of size two (for n = 3k + 1) or three (for n = 3k + 2). In the first case, we can eliminate two old triangles at the cost of creating four new ones; hence, the new arrangement \mathcal{A}_{3k+1} has n = 3k + 1 pseudocircles and $4k + 2 = \lceil \frac{4}{3}n \rceil$ triangles. In the second case, we can eliminate three old triangles at the cost of creating six new ones and obtain a new arrangement \mathcal{A}_{3k+2} with n = 3k + 2 pseudocircles and $4k + 3 = \lceil \frac{4}{3}n \rceil$ triangles.

4.2 Proof of Proposition 7

In this subsection we construct intersecting digon-free arrangements with few triangles (less than 2n-4) and no \mathcal{N}_6^{Δ} subarrangement. The key to the construction is again the replacement of pseudocircles of a base arrangement by bundles. In the following proof we use bundles of size 3. In the discussion we will also mention size 4 and larger sizes.

▶ Proposition 7. There is an infinite family of simple intersecting digon-free arrangements of n pseudocircles with $\lceil \frac{5}{3}n \rceil + 2$ triangles which have no subarrangement isomorphic to \mathcal{N}_6^{Δ} .

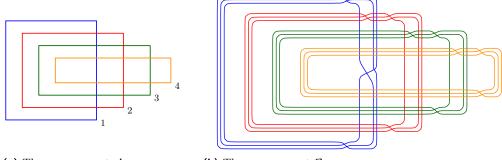
Proof. For any $N \geq 2$, let A_N be the arrangement of pseudocircles C_1, \dots, C_N such that the cyclic order of intersection of C_i with the other pseudocircles is:

$$1, 2, \ldots, i-1, i+1, \ldots, N, N, N-1, \ldots, i+1, i-1, \ldots, 1.$$

The arrangement can be represented with axis-parallel rectangles so that for all i < j the right side of C_i cuts vertically through C_j ; see Figure 18a for \mathcal{A}_4 . Note that every triple of

three pseudocircles of A_N induces a NonKrupp. The faces of A_N are: 3 digons, 2(N-2) triangles, N(N-3) + 2 four-gons, and a single (2N-2)-gon.

On the basis of A_N we construct an arrangement B_N by bundle replacements.



(a) The arrangement A_4

(b) The arrangement \mathcal{B}_4 .

Figure 18 Replacing the pseudocircles with bundles so that the twists destroy all digons.

From \mathcal{A}_N we obtain the arrangement \mathcal{B}_N by successively replacing C_1, \dots, C_N by bundles B_1, \dots, B_N of size 3 where we place the twists according to the following rules; Figure 18b shows the arrangement \mathcal{B}_4 obtained by using the base arrangement \mathcal{A}_4 from Figure 18a.

- The bundle replacing C_1 , has two twists outside of C_2 and one twist inside of C_N .
- The bundle B_i replacing C_i , 1 < i < N, has two twists which increase the degree of each of the two triangles formed by B_{i-1} , C_i and C_{i+1} . The third twist is within B_{i-1} such that it increases the degree of a triangle corresponding to a twist of B_{i-1} .
- Finally, when replacing C_N , place two consecutive twist which increase the degree of the digon formed by B_{N-1} and C_N and one twist in a triangle corresponding to a twist of B_{N-1} .

The arrangement \mathcal{B}_N consists of n=3N pseudocircles. For counting the triangles, each bundle B_i has 6 internal triangles, but for the bundles B_1, \dots, B_{N-1} one of them is dissolved by a twist of the next bundle B_{i+1} . Additionally, the digon formed by C_1 and C_N has been made a triangle. The total number of triangles of \mathcal{B}_N is

$$6N - (N - 1) + 1 = 5N + 2 = \frac{5}{3}n + 2.$$

The arrangement \mathcal{B}_N is digon-free. It remains to show that there is no subarrangement of \mathcal{B}_N which is isomorphic to \mathcal{N}_6^{Δ} .

Since every triple of pseudocircles of \mathcal{A}_N induces a NonKrupp, the same is true for triples of pseudocircles taken from distinct bundles of \mathcal{B}_N . A \mathcal{N}_6^{Δ} arrangement has exactly 4 triples that form a NonKrupp and each of the six pseudocircles is member of exactly two of them.

Let \mathcal{B}' be a subarrangement of \mathcal{B}_N consisting of 6 pseudocircles. If \mathcal{B}' contains pseudocircles of at least four different bundles then it has strictly more than 4 NonKrupps, hence, it cannot be isomorphic to \mathcal{N}_6^{Δ} . Now assume that \mathcal{B}' consists of $k_1, k_2, k_3 \geq 0$ pseudocircles from three pairwise different bundles, $k_1 + k_2 + k_3 = 6$.

 \triangleright Claim 14. If $k_i > 0$ for i = 1, 2, 3, then \mathcal{B}' is not isomorphic to \mathcal{N}_6^{Δ} .

Proof of Claim 14. As each triple of pseudocircles of pairwise different bundles forms a NonKrupp, \mathcal{B}' contains at least $k_1 \cdot k_2 \cdot k_3$ NonKrupp subarrangements. Since $k_1 + k_2 + k_3 = 6$, this value is 4, 6, or 8. If $k_1 \cdot k_2 \cdot k_3 = 4$, two of the k_i are equal to 1 and the two corresponding pseudocircles participate in all 4 NonKrupp subarrangements of \mathcal{B}' . Hence, in all cases \mathcal{B}' and \mathcal{N}_6^{Δ} are not isomorphic.

It follows that if \mathcal{B}' is isomorphic to \mathcal{N}_6^{Δ} , then it must contain the three pseudocircles of each of two bundles, i.e., $k_1 = k_2 = 3$.

 \triangleright Claim 15. If \mathcal{B}' is a subarrangement consisting of two complete bundles of \mathcal{B}_N , then \mathcal{B}' is not isomorphic to \mathcal{N}_6^{Δ} .

Proof of Claim 15. From the bundle structure of \mathcal{B}' we get 6 triangles in each of the two bundles. A twist of one bundle placed between consecutive pseudocircles of the other bundle can destroy one triangle of the second bundle. Such a twist corresponds to one of the two crossings in the underlying arrangement of two circles. Hence the arrangement \mathcal{B}' has at least 10 triangles, while \mathcal{N}_6^{Δ} has only 8.

This excludes the existence of a subarrangement \mathcal{B}' isomorphic to \mathcal{N}_6^{Δ} .

We now sketch how the constant 5/3 of Proposition 7 could be replaced by the smaller constant 3/2.

Again we take the arrangement A_N as a basis but now we replace each pseudocircle with a bundle of size 4. This can lead to an intersecting digon-free arrangement with n=4N pseudocircles and

$$(8N + 2(N - 2) + 6) - 4N = 6N + 2 = \frac{3}{2}n + 2$$

triangles; the count is as follows: each of the 4N twists can increase the degree of a triangle or digon by one, the initial arrangement \mathcal{A}_N has 3 digons and 2(N-2) triangles. To achieve \mathcal{N}_6^{Δ} -freeness with bundles of size 4 the twists have to be placed with some care, and the analysis that the result is indeed \mathcal{N}_6^{Δ} -free requires the analysis of a lot of cases. In fact \mathcal{N}_6^{Δ} can be obtained from the arrangement with two circles by replacing one with a bundle of size 2 and the other with a bundle of size 4 if the twists are placed in a specific way.

Every arrangement \mathcal{A} with the property that every triple of pseudocircles forms a NonKrupp can be used as the basis for a construction with bundle replacement with bundles of size 3 and/or 4 such that the constructed arrangement is intersecting, digon-free, \mathcal{N}_6^{Δ} -free and has few triangles. Using bundles of larger size makes it more challenging to avoid \mathcal{N}_6^{Δ} -subarrangements, and we see no way of getting below $\frac{3}{2}n$ triangles with such a construction.

For Conjecture 6 to be true, it would be necessary that \mathcal{N}_6^{Δ} -free arrangements obtained by bundle replacement with few triangles are non-circularizable. With the help of the polymake [9] extension r9n developed by Julian Pfeifle, we could verify that the arrangement \mathcal{B}_3 with 9 pseudocircles and 17 triangles and the arrangement \mathcal{C}_7 with 7 pseudocircles shown in Figure 19 are both not circularizable. We leave the following questions for future research:

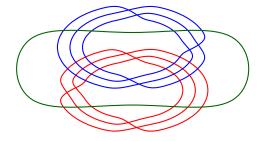


Figure 19 Non-circularizable arrangement C_7 .

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- What is the minimum number of triangles of intersecting, digon-free, \mathcal{N}_6^{Δ} -free arrangements of pseudocircles?
- What is the minimum number of triangles of intersecting, digon-free arrangements of circles? Is it 2n-4? (Conjecture 6)

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